

VARIABLE LENGTH DECODING SYSTEM AND METHOD

CROSS-REFERENCE TO RELATED APPLICATIONS

This application claims priority from provisional application No. 60/462,501, filed 04/11/2003 and is a continuation-in-part of pending application No. 09/788,807, filed 02/20/2001.

BACKGROUND OF THE INVENTION

The present invention relates to electronic systems, and more particularly, to digital systems and methods with bitstreams representing coded information with codewords of variable length.

The current rapid expansion of digital communication (speech, video, and data) relies on increasingly economical digital signal processing and efficient transmission and storage. For example, video communication has general functionality as illustrated in Figure 4a, and increasingly includes a link through the air interface as illustrated in Figure 4b. Many digital communication systems and standards, such as MPEG, use coding with variable length codewords for coding efficiency. Variable length decoding (VLD) is needed for decoding bitstreams whenever variable length coding (VLC) is used by the encoder for generating the bitstreams. A VLC table typically has entries with three fields (codeword length: *length*; pattern or information encoded: *pattern*; and variable length codeword: *vlc_code*). VLD is to determine the value of the fields (*length*, *pattern*) based on the *vlc_code* value extracted from the bitstream. Figure 2 illustrates the principle of VLD. To find and decode the next codeword, the decoder looks at the sequence of bits forward from the current bitstream position and finds a match to a possible value of *vlc_code*. Based on the extracted *vlc_code* value, the VLD determines (*length*, *pattern*) by look up in the VLC table, outputs the value of *pattern* as the decoded next codeword, and then updates the current decoding position in the bitstream according to the decoded codeword *length*, and starts to decode the next codeword. The look for the next codeword usually reads a fixed

length of bits (e.g., the maximal codeword length, *len_max*, of the entire VLC table) and then searches for a possible codeword; see Figure 2.

Normally, a decoder includes several VLDs to handle multiple VLCs because VLC tables are different from table to table, and VLD functions have to be implemented differently according to the contents of the VLC tables. For a decoder implementation this implies large code size (or high gate count for hardware solutions) and long development times. Therefore, there is a demand to have a universal VLD method that is able to deal with any VLC table in order to reduce costs and increase flexibility in decoder design.

Obviously, the simplest way to do universal VLD is with the direct VLD table look up; that is, each possible sequence of *len_max* bits in the bitstream is an index to a table entry containing the next codeword *length* and *pattern*. However, this requires a huge VLD table size: indeed, a table with 2^{len_max} entries. For example, if the maximum codeword size of a VLC table is 16 bits (i.e., *len_max* = 16), such a VLD table would have 64 K entries. This is too expensive in terms of memory size.

SUMMARY OF THE INVENTION

The present invention provides universal VLD methods including a VLD table construction function and a universal VLD function. The universal VLD function is valid for any VLD as long as the VLD tables are produced by using the VLD table construction function.

This has the advantage of smaller VLD size because a single VLD function can decode multiple VLC codes.

BRIEF DESCRIPTION OF THE DRAWINGS

Figures 1a-1c illustrate preferred embodiment universal variable length decoding.

Figure 2 illustrates variable length decoding (VLD).

Figures 3a-3b are block diagrams of a preferred embodiment decoding systems.

Figures 4a-4b show general digital communication which could use preferred embodiment decoding.

DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

1. Overview

Preferred embodiments include universal variable length decoders and universal variable length decoding (VLD) methods for decoding streams of codewords encoded with variable length coding (VLC). A given VLC table is translated into a VLD control table plus a VLD code table, and a universal variable length decoding (UVLD) use these VLD tables to decode a stream of VLC codewords; see Figure 1a. A single decoder for bitstreams with differing VLC table encodings simplifies the decoding.

For a given VLC table which is either a 0-leading code table or a 1-leading code table, the VLD code table is not too large and UVLD directly applies. However, for some VLC tables the corresponding VLD code table may require large memory. For such VLC tables, advanced UVLD (AUVLD) partitions a VLC table into prefix-oriented tables with a VLD prefix table plus VLD control and code table construction applied to each prefix-oriented table; see Figure 1b. AUVLD can effectively handle VLC tables such as the RVLC of DCT coefficients in MPEG-4. And the VLD prefix table can be merged with the VLD control table to minimize storage requirements; see Figure 1c.

For UVLD the bitstream provides addressing into the VLD control table entries (*shift, offset*), and these provide addressing into the VLD code table entries (*length, pattern*) for decoding and moving the read position in the bitstream. Additionally, a *reverse* indicates bit complementation from 1-leading to 0-leading.

The AUVLD adds prefix parameter (*pbits*) indicates the number of bits in the bit pattern, *prefix*, which defines the prefix-oriented tables of a partitioned VLC table, and a reorganization mask (*mask_key*) provides further memory reduction by optimizing the sizes of the VLD tables constructed from a prefix-oriented table.

The following first considers UVLD without prefixes and then describes AUVLD which extends UVLD with the prefix partitioning of a VLC table.

2. Construction of VLD code table and VLD control table

Basically, there are two kinds of simple VLC tables: 0-leading tables (e.g. Table 1) and 1-leading tables (e.g. Table 4). First consider the 0-leading tables; a bit reversal will convert 1-leading tables into 0-leading tables. Following sections will consider extensions to more general prefix-oriented VLC tables such as Table 8.

Now an example of a 0-leading VLC table is the following Table 1 which is the MPEG-1 macroblock address increment code table.

length	pattern (macroblock_address_increment)	vlc_code (codeword)
1	1	1
3	2	011
3	3	010
4	4	0011
4	5	0010
5	6	00011
5	7	00010
7	8	0000111
7	9	0000110
8	10	00001011
8	11	00001010
8	12	00001001
8	13	00001000
8	14	00000111
8	15	00000110
10	16	0000010111
10	17	0000010110
10	18	0000010101
10	19	0000010100
10	20	0000010011
10	21	0000010010
11	22	00000100011
11	23	00000100010
11	24	00000100001
11	25	00000100000
11	26	00000011111
11	27	00000011110
11	28	00000011101
11	29	00000011100
11	30	00000011011
11	31	00000011010
11	32	00000011001
11	33	00000011000

Table 1. VLC table of MPEG1 Macroblock_address_increment

Table 1 is used as an example to explain how to construct the VLD code table and the VLD control table from a given VLC table.

The basic idea is to divide the VLC table into a set of sub-tables in a defined order according to the *vlc_code* values. The sub-table entries make up the VLD code table. For each sub-table there will be a field to indicate the location of the sub-table in the VLD code table, and those fields build up the VLD control table.

The construction of the VLD code table and the VLD control table from a given VLC table has eight steps:

1. Get the maximal length (*len_max*) of codewords in the VLC table (e.g., *len_max* = 11 for Table 1); and let *length* denote the length of a codeword.
2. Left shift each VLC codeword (*vlc_code*) in the VLC table by (*len_max* – *length*) bits. For example, the first *vlc_code* entry in Table 1 is 1 and has a length of 1 bit, so the codeword is left-shifted by 10 bits to yield 1--- ---- ---. After shifting this codeword has a value of 1024 when interpreted as an 11-bit integer; see the last entry in columns "Re-organized *vlc_code*" and "shifted value" in following Table 2 which illustrates items in the construction of the VLD tables from Table 1. That is, treat the left shifting as adding 10 0-bits, so the shifted codeword would be 1000 0000 000 which, as a binary integer, equals 1024.
3. Reorder the shifted VLC codewords into increasing order according to the *vlc_code* values after shifting (column "shifted value" in Table 2).
4. Divide the VLC table into sub-tables according to the shifted values. A reorganized VLC codeword is classified into sub-table_{*n*} if its shifted value satisfies $2^{n-1} \leq \text{value} < 2^n$. The variable *subtab_id* is used to denote *n* in the decoding. Sub-table₀ is an exception and contains only one zero element. By classifying sub-tables in this way it is easy to identify to which sub-table index, *subtab_id*, a given shifted value belongs, because the sub-table index can be simply determined by checking the MSB position of the shifted value. This is the same as the sub-table number for shifted *vlc_code* being 11 – (the number of leading 0s in *vlc_code*).

5. Fill up the leaks in each sub-table. Leaks are defined as codeword entries that are valid but not used in the VLC table. For example, in Table 2 all the entries in sub-table0, sub-table1, sub-table2, sub-table3, and sub-table4 plus the entries below shifted value 24 in sub-table5 are missing because those entries are not used in the VLC table (see Table 1). For the purpose of error detection, those leaks must be filled with dummy codewords. In Table 2, the leaks are filled with the dummy code (0,0) for $(length, pattern)$ in the central column and correspond to potential-but-not-used codewords shown underlined in column “Re-organized vlc_code”. Whenever a decoder extracts a dummy code (0,0) from the bitstream, it will report a finding of error.
6. Determine $shift$ for each sub-table where $shift$ for a sub-table is defined as the difference between len_max and the maximal length of the codewords in the sub-table. For example, sub-table6 of Table 2 has codewords varying in length from 8 to 11 bits, thus $shift$ equals 0. The VLD control table has $shift$ as the first field component; see Table 2, left column “VLD-control (shift, offset)”. In effect, $shift$ is the minimum shift of all of the codewords in the sub-table and will be applied in the decoding to convert a set of bits of length len_max from the bitstream to approximate codeword size. For example, sub-table7 has vlc_code varying from 7 to 8 bits, so $shift = 3$; then for $vlc_code = 0000\ 111$ the corresponding len_max bits from the bitstream ($0000\ 111x\ xxx$) after left shifting by $shift$ would be $0000\ 111x$ where x is either 0 or 1.
7. Determine the entry indices in the VLD code table by simple enumeration starting at 0 in sub-table0. Each element $(length, pattern)$ is repeated $2^{len_max - shift - length}$ times in the VLD code table; this means the approximate codeword suffices because the repetition allows for the irrelevant x . As an example, in sub-table7 of Table 2, the center column (which are the VLD code table entries $(length, pattern)$ expressed in decimal) and Table 3 which lists these entries, the entry (7, 8) appears twice because ($len_max = 11, shift = 3, length = 7$). That is, the VLD code table indices 70 and 71

(column "VLD code index" of Table 2) both appear for (7,8). and in the corresponding VLD code table (VLDCodeTab[79] of Table 3), the 70th and 71st components in the array of 79 components are both (7,8). This repetition accounts for the unequal length of codewords in the sub-table (the irrelevant x bits) and makes universal decoding simpler.

8. Determine *offset* for each sub-table so that any entry in the sub-table can find its *index* in the VLD code table by the equation: $index = offset + (value \gg shift)$, where *value* is the entry in the column "shifted value" in Table 2. Thus *offset* for a sub-table aligns the *index* with the unshifted *value*. Note that decoding "shifted value" here is equivalent to the value of the sequence of *len_max* bits starting at the current decoding position in the bitstream in that the binary value of the *len_max* bits is in the range of the sub-table. For example, the second entry of sub-table10 of Table 2 has shift 8, index 77, and *vlc_code* 011- ---- -; the 8 -'s are the left shift of 8, so ($value \gg shift$) is 3 (binary 011) and $offset = index - (value \gg shift) = 77 - 3 = 74$. The VLD control table has *offset* as the second field component.

VLD control (shift,offset)	VLD code index	VLD code (length,pattern)	Re-organized vlc_code	shifted value
Sub-table0 (0~0)				
(0,0)	0~0	(0, 0)	0000 0000 000	0
Sub-table1 (1~1)				
(0,0)	1~1	(0, 0)	0000 0000 001	1
Sub-table2 (2~3)				
(0,0)	2~3	(0, 0)	0000 0000 010	2~3
		(0, 0)	0000 0000 011	
Sub-table3 (4~7)				
(0,0)	4~7	(0, 0)	0000 0000 100	4~7
		~	~	
		(0, 0)	0000 0000 111	
Sub-table4 (8~15)				
(0,0)	8~15	(0, 0)	0000 0001 000	8~15
		~	~	
		(0, 0)	0000 0001 111	
Sub-table5 (16~31)				
(0,0)	16~23	(0, 0)	0000 0010 000	16~23
		~	~	
		(0, 0)	000 0010 111	
	24	(11,33)	0000 0011 000	24
	25	(11,32)	0000 0011 001	25
	26	(11,31)	0000 0011 010	26
	27	(11,30)	0000 0011 011	27

	28	(11, 29)	0000 0011 100	28
	29	(11, 28)	0000 0011 101	29
	30	(11, 27)	0000 0011 110	30
	31	(11, 26)	0000 0011 111	31
Sub-table6 (32~63)				
(0, 0)	32	(11, 25)	0000 0100 000	32
	33	(11, 24)	0000 0100 001	33
	34	(11, 23)	0000 0100 010	34
	35	(11, 22)	0000 0100 011	35
	36~37	(10, 21)	0000 0100 10-	36
	38~39	(10, 20)	0000 0100 11-	38
	40~41	(10, 19)	0000 0101 00-	40
	42~43	(10, 18)	0000 0101 01-	42
	44~45	(10, 17)	0000 0101 10-	44
	46~47	(10, 16)	0000 0101 11-	46
	48~55	(8, 15)	0000 0110 ---	48
	56~63	(8, 14)	0000 0111 ---	56
Sub-table7 (64~127)				
(3, 56)	64	(8, 13)	0000 1000 ---	64
	65	(8, 12)	0000 1001 ---	72
	66	(8, 11)	0000 1010 ---	80
	67	(8, 10)	0000 1011 ---	88
	68~69	(7, 9)	0000 110- ---	96
	70~71	(7, 8)	0000 111- ---	112
Sub-table8 (128~255)				
(6, 70)	72	(5, 7)	0001 0--- ---	128
	73	(5, 6)	0001 1--- ---	192
Sub-table9 (256~511)				
(7, 72)	74	(4, 5)	0010 ---- ---	256
	75	(4, 4)	0011 ---- ---	384
Sub-table10 (512~1023)				
(8, 74)	76	(3, 3)	010- ---- ---	512
	77	(3, 2)	011- ---- ---	768
Sub-table11 (1024~2047)				
(10, 77)	78	(1, 1)	1--- ---- ---	1024

**Table 2. Design of UVLD code table and VLD control table for MPEG1
Macroblock_address_increment**

Given the following data types:

```
typedef struct vldcodetab {
    char length;
    short pattern;
} VLDCodeTab;

typedef struct vldctltab {
    char shift;
    short offset;
} VLDCtlTab;
```

VLD code table and VLD control table for VLC Table 1 are shown in Table 3.

```
static VLDCodeTab Macroblock address_increment vldtab[79]={
```

```

    { 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},
    { 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},
    { 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},
    {11,33},{11,32},{11,31},{11,30},{11,29},{11,28},{11,27},{11,26},
    {11,25},{11,24},{11,23},{11,22},{10,21},{10,21},{10,20},{10,20},
    {10,19},{10,19},{10,18},{10,18},{10,17},{10,17},{10,16},{10,16},
    { 8,15},{ 8,15},{ 8,15},{ 8,15},{ 8,15},{ 8,15},{ 8,15},{ 8,15},
    { 8,14},{ 8,14},{ 8,14},{ 8,14},{ 8,14},{ 8,14},{ 8,14},{ 8,14},
    { 8,13},{ 8,12},{ 8,11},{ 8,10},{ 7, 9},{ 7, 9},{ 7, 8},{ 7, 8},
    { 5, 7},{ 5, 6},{ 4, 5},{ 4, 4},{ 3, 3},{ 3, 2},{ 1, 1}
};

static VLDCtlTab Macroblock_address_increment_vldctl[12]={
    { 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 0, 0},{ 3,56},
    { 6,70},{ 7,72},{ 8,74},{10,77}
};

len_max = 11; reverse = 0;

```

Table 3. UVLD code table and UVLD control table for MPEG1 macroblock_address_increment

That is, the VLD construction function converts VLC Table 1 into the VLD code table and the VLD control table shown in Table 3. Note that the original VLC Table 1 has 33 entries (*length, pattern, vlc_code*); whereas, the VLD code table has 79 entries (*length, pattern*) and the VLD control table has 12 entries (*shift, offset*). In effect, the preferred embodiment has translated VLC Table 1 into two tables which may aggregately be larger, but this translation will allow application of a universal decoding function. Note that the number of repetitions in the VLD code table depends upon the variance of codeword length for codewords with the same number of leading 0s (which will be in the same sub-table). Section 3 describes a universal VLD decoding using the VLD control table and the VLD code table instead of a decoder specialized for the original VLC table.

Now consider the VLD construction function applied to a 1-leading VLC table such as the following Table 4 example which is a 1-leading VLC table and used for the MPEG1 dct_dc_size_luminance.

length	pattern (dct_dc_size_luminance)	vlc_code
3	0	100
2	1	00
2	2	01

3	3	101
3	4	110
4	5	1110
5	6	11110
6	7	111110
7	8	1111110

Table 4. VLC table of MPEG1 dct_dc_size_luminance

The VLD table construction for a 1-leading VLC table is similar to that of a 0-leading VLC table. The difference lies in re-organization of the *vlc_codes*. In a 1-leading table, the shifted value of a *vlc_code* is computed by first performing a bit-reversal (interchange 0 and 1 and indicated by the value of *reverse*) of the *vlc_code* followed by a left shift of (*len_max* – *length*) bits (see Table 5, column "shifted value"), where *length* is again the codeword length of *vlc_code*. After this computation of shifted value, the same steps as described above for 0-leading codes are used to construct the VLD code table and the VLD control table: see Table 5 and Table 6.

VLD control (shift,offset)	VLD code index	VLD code (length,pattern)	Re-organized vlc_code	Shifted Value
Sub-table0 (0~0)				
(0,0)	0	(0, 0)	1111 111	0
Sub-table1 (1~1)				
(0,0)	1	(7, 8)	1111 110	1
Sub-table2 (2~3)				
(1,1)	2	(6, 7)	1111 10-	2
Sub-table3 (4~7)				
(2,2)	3	(5, 6)	1111 0--	4
Sub-table4 (8~15)				
(3,3)	4	(4, 5)	1110 ---	8
Sub-table5 (16~31)				
(4,4)	5	(3, 4)	110- ---	16
Sub-table6 (32~63)				
(4,4)	6	(3, 3)	101- ---	32
	7	(3, 0)	100- ---	48
Sub-table7 (64~127)				
(5,6)	8	(2, 2)	01-- ---	64
	9	(2, 1)	00-- ---	96

Table 5. Design of VLD code table and VLD control table for MPEG1 dct_dc_size_luminance

```
static VLDCodeTab Dct_dc_size_luminance_dcdtab[10]={
    { 0, 0},{ 7, 8},{ 6, 7},{ 5, 6},{ 4, 5},{ 3, 4},{ 3, 3},{ 3, 0},
    { 2, 2},{ 2, 1}
```

```

};

static VLDCtlTab Dct_dc_size_luminance_dcdctl[8]={
    { 0, 0},{ 0, 0},{ 1, 1},{ 2, 2},{ 3, 3},{ 4, 4},{ 4, 4},{ 5, 6}
};

len_max = 7; reverse = 1;

```

Table 6. VLD Code table and VLD control table for MPEG1 dct_dc_size_luminance

As Table 3 and Table 6 show, the size of a VLD control table is fixed; it has $len_max + 1$ entries. However, the size of a VLD code table varies from table to table, it depends on the VLC table characteristics.

The construction of the VLD code table and the VLD control table as in the foregoing for any given VLC table can be automated; this includes the decision whether the input VLC table is treated as a 0-leading table or a 1-leading table.

3. Universal VLD decoding

The construction of the VLD code table and the VLD control table for a VLC table allows a universal VLD decoding function for decoding. Table 7 illustrates the pseudo code for an implementation of a universal VLD decoding function.

```

int UniversalVLD(
    Bitstream *stream,          /* pointer of bitstream */
    VLDCodeTab *vldtab,        /* pointer of VLD table */
    VLDCtlTab *vldctl,         /* pointer of VLD control table */
    int len_max,                /* maximum code length in the VLD table */
    char reverse,               /* reverse =0/1 -> zero/one leading VLD table */
    char *err_flag)             /* err_flag =1 ->error detected, err_flag=0->decoding OK */
{
    int value, subtab_id, index;
    /*=====*/
    /* get the value of next "len_max" bits in the bitstream */
    /*=====*/
    value = next_bits(stream, len_max);

    /*=====*/
    /* reverse the value for 1-leading VLD table */
    /*=====*/
    if (reverse) value = ((1<<len_max)-1)^value;

    /*=====*/
    /* determine the sub-table index according to the value. TMS320C6X and TMS320C54X have */
    /*=====*/
}

```

```

/* special instructions for such an operation */
/*=====*/
if (value==0) subtab_id = 0; else subtab_id = (int) log2(value) + 1;

/*=====*/
/* get index in the VLD code table */
/*=====*/
index = vldctl[subtab_id].offset + (value>>vldctl[subtab_id].shift);

/*=====*/
/* decide if an decoding error is detected */
/*=====*/
if (vldtab[index].length ==0) *err_flag =1; else *err_flag=0;

/*=====*/
/* update the current decoding position in the bitstream */
/*=====*/
if (*err_flg==0) flush_bits(stream, vldtab[index].length);

/*=====*/
/* return the decoded coding pattern */
/*=====*/
return vldtab[index].pattern;
}

```

Table 7. Pseudo code for the universal VLD decoding function

The universal VLD decoding function contains the following steps:

1. Get the *value* of the next *len_max* bits in the bitstream; that is, interpret the next *len_max* bits as a binary integer.
2. If the VLC is a 1-leading table, then reverse the bits of *value* bit-by-bit.
3. Determine the sub-table index (*subtab_id*) for addressing the VLD control table according to the *value*: if *value* = 0, then the *subtab_id* is 0; otherwise, $subtab_id = (int) \log_2(value) + 1$.
4. Get *shift* and *offset* from the VLD control table by using the *subtab_id* address.
5. Compute the *index* in the VLD code table as $index = (value \gg shift) + offset$.
6. Acquire the codeword *length* and codeword *pattern* from the VLD code table entry at *index*.
7. Update the current decoding position in the bitstream by using *length*, and interpret *pattern* to recover the encoded symbol.
8. Loop to step 1.

This function is universal in the sense that it can deal with any VLC table provided that corresponding VLD code table and VLD control table are constructed according to the foregoing description. Further, this function has the ability to perform error detection. Figure 1a illustrates the decoding.

Figure 3a is a block diagram of preferred embodiment universal VLD decoding function. The core of this function is the universal VLD unit, which includes data memory and registers connected to it. The data memory is used to store the VLD code table and the VLD control table, while the three registers are used to store the bitstream decoding position, *len_max/reverse*, and the decoding status (i.e. *err_flag*). To decode *pattern*, the universal VLD decoding unit points to the related VLD code table and VLD control table, then extracts the *pattern* from the bitstream according to the given bitstream position and *len_max/reverse*. After decoding the *pattern*, the bitstream position register as well as the decoding status register are updated.

The preferred embodiment universal VLD system is made up of two functions: a VLD table construction function and a universal VLD decoding function. The VLD table construction function constructs the VLD code table and the VLD control table according to the given VLC table; this can be done offline and stored prior to actual decoding. The universal VLD decoding function is valid for decoding any VLC as long as its VLC table is translated into a VLD code table plus a VLD control table according to the foregoing format. In addition, it provides the error detection ability that is essential for decoding a VLC in real applications.

4. Advanced UVLD

Further preferred embodiment universal variable length decoding has a two-tier decomposition of a VLC table: each codeword is assigned to a prefix-oriented (prefix-labeled) table defined by a fixed-bit-length prefix (left-most bits) of the codeword. Then from each prefix-oriented table construct a control table and a code table as described in the preceding UVLD sections. Also, an auxiliary table of prefixes and mask keys (used for reorganizations of codewords

within prefix-oriented tables prior to control table and code table constructions) provides the addressing to the appropriate prefix-oriented table constructs from the bitstream. This extends the UVLD described in the preceding sections for 0-leading or 1-leading VLC tables which can be considered as having a single prefix-oriented table with a prefix length of 0 and a mask key of all 0s or all 1s according to *reverse*, respectively. The preferred embodiments are termed Advanced UVLD or AUVLD and overcome a decoder memory size problem with UVLD for certain VLC codes by computing the required tables for each possible prefix length and picking the prefix length which minimizes memory use.

In particular, presume a given VLC table, then for each value of the integer *pbits* in the range $0 \leq pbits < len_max$, where *len_max* is the maximum codeword bit length (exclusive of a suffixed sign bit) in the VLC table, proceed as follows:

1. Divide the given VLC table into 2^{pbits} prefix-oriented tables by using the first (left) *pbits* bits of a codeword (exclusive of a suffixed sign bit) as its *prefix*; strip off the *prefix* bits to define the codeword in the prefix-oriented table.
2. Within each prefix-oriented table, apply a *mask_key* to reorganize the table entries; repeat for all possible mask keys with selection of *mask_key* to minimize memory requirements.
3. After the prefix-oriented table reorganization based on *mask_key*, use the prefix-oriented table to construct a VLD code table and a VLD control table as previously described for a general VLC table; that is, an AUVLD control table and an AUVLD code table are constructed for each prefix-oriented table with each *mask_key*.
4. Compute memory required for storage of the various tables and pick the values of *pbits* and available *mask_key* which minimize the memory requirement; this defines the AUVLD control and code tables for the prefix-oriented code tables. Also include a AUVLD prefix table which contains the *mask_key* and translates bitstream bits into the addresses to entries in the control and code tables plus *mask_key* for decoding.

In more detail, for a given value for $pbits$ plus a given VLC table, first decompose the VLC table into 2^{pbits} prefix-oriented tables. Let $(length, pattern, vlc_code)$ represent an entry in the VLC table where $length$ is the codeword length in bits (excluding any sign-bit suffix), $pattern$ is the symbol being encoded (such as the last/run/level of Table 8), and vlc_code is the corresponding codeword; also let len_max be the maximum of $length$ in the VLC code table. A VLC code table entry $(length, pattern, vlc_code)$ is classified into the prefix-oriented table labeled $prefix$ when the first $pbits$ bits of the codeword (with left padded 0s if needed) have a value equal to $prefix$; that is, $(vlc_code >> (length - pbits)) = prefix$.

A prefix-oriented table can have none, one, or multiple entries, depending upon the given VLC code table. The first $pbits$ of each codeword are then stripped off to form the entry in the prefix-oriented table. Therefore, the counterpart of an entry $(length, pattern, vlc_code)$ from the VLC code table is the entry $(length-pbits, pattern, vlc_code \& (2^{length-pbits} - 1))$ in the appropriate prefix-oriented table where $\&$ indicates bitwise AND.

For a more detailed explanation, consider the example of the MPEG-4 reversible variable length coding of the DCT coefficients for intercoded frame macroblocks; the VLC code table (Table 8) has 170 entries. This is a run-length type of coding where a DCT coefficient to be encoded is represented by three numbers: “last” (1-bit), “run” (6-bit), and “level” (5-bit) which make up the 12-bit $pattern$ of the entry in the VLC code table. Disregarding any suffixed sign bit, $length$ varies from 3 to 15; see following Table 8. Note that in the table the 12-bit pattern and the codeword are both expressed in decimal for compactness.

Static CodTab MPEG4_RVLC_INTER_DCT_TAB[170]={ /* {code_length, code_word, code_pattern}

```
{ 3, 6, 1}, { 3, 7, 33}, { 4, 1, 2}, { 4, 10, 65}, { 4, 11,2049},
{ 5, 4, 3}, { 5, 5, 97}, { 5, 8, 129}, { 5, 9, 161}, { 5, 18,2081},
{ 5, 19,2113}, { 6, 12, 34}, { 6, 13, 193}, { 6, 20, 225}, { 6, 21, 257},
{ 6, 24,2145}, { 6, 25,2177}, { 6, 34,2209}, { 6, 35,2241}, { 7, 28, 4},
{ 7, 29, 66}, { 7, 44, 289}, { 7, 45, 321}, { 7, 52, 353}, { 7, 53,2273},
{ 7, 56,2305}, { 7, 57,2337}, { 7, 66,2369}, { 7, 67,2401}, { 8, 60, 5},
{ 8, 61, 6}, { 8, 92, 35}, { 8, 93, 98}, { 8, 108, 130}, { 8, 109, 385},
{ 8, 116, 417}, { 8, 117, 449}, { 8, 120,2050}, { 8, 121,2433}, { 8, 130,2465},
{ 8, 131,2497}, { 9, 124, 7}, { 9, 125, 36}, { 9, 188, 67}, { 9, 189, 162},
{ 9, 220, 481}, { 9, 221, 513}, { 9, 236, 545}, { 9, 237,2082}, { 9, 244,2529},
{ 9, 245,2561}, { 9, 248,2593}, { 9, 249,2625}, { 9, 258,2657}, { 9, 259,2689},
{10, 252, 8}, {10, 253, 9}, {10, 380, 37}, {10, 381, 99}, {10, 444, 194},
```



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{10, 445, 226}, {10, 476, 258}, {10, 477, 290}, {10, 492, 577}, {10, 493, 609},
{10, 500, 641}, {10, 501,2721}, {10, 504,2753}, {10, 505,2785}, {10, 514,2817},
{10, 515,2849}, {11, 508, 10}, {11, 509, 11}, {11, 764, 38}, {11, 765, 68},
{11, 892, 131}, {11, 893, 163}, {11, 956, 322}, {11, 957, 673}, {11, 988, 705},
{11, 989, 737}, {11, 1004, 769}, {11, 1005, 801}, {11, 1012, 833}, {11, 1013,2051},
{11, 1016,2114}, {11, 1017,2881}, {11, 1026,2913}, {11, 1027,2945}, {12, 1020, 12},
{12, 1021, 39}, {12, 1532, 69}, {12, 1533, 100}, {12, 1788, 195}, {12, 1789, 227},
{12, 1916, 354}, {12, 1917, 865}, {12, 1980, 897}, {12, 1981, 929}, {12, 2012,2083},
{12, 2013,2146}, {12, 2028,2178}, {12, 2029,2977}, {12, 2036,3009}, {12, 2037,3041},
{12, 2040,3073}, {12, 2041,3105}, {12, 2050,3137}, {12, 2051,3169}, {13, 2044, 13},
{13, 2045, 14}, {13, 3068, 15}, {13, 3069, 16}, {13, 3580, 40}, {13, 3581, 101},
{13, 3836, 132}, {13, 3837, 164}, {13, 3964, 259}, {13, 3965, 386}, {13, 4028, 961},
{13, 4029, 993}, {13, 4060,1025}, {13, 4061,1057}, {13, 4076,2052}, {13, 4077,2210},
{13, 4084,2242}, {13, 4085,2274}, {13, 4088,2306}, {13, 4089,2338}, {13, 4098,3201},
{13, 4099,3233}, {14, 4092, 17}, {14, 4093, 18}, {14, 6140, 41}, {14, 6141, 42},
{14, 7164, 70}, {14, 7165, 71}, {14, 7676, 102}, {14, 7677, 196}, {14, 7932, 291},
{14, 7933, 418}, {14, 8060, 450}, {14, 8061, 482}, {14, 8124, 514}, {14, 8125,1089},
{14, 8156,1121}, {14, 8157,1153}, {14, 8172,2053}, {14, 8173,2084}, {14, 8180,2370},
{14, 8181,2402}, {14, 8184,2434}, {14, 8185,3265}, {14, 8194,3297}, {14, 8195,3329},
{15, 8188, 19}, {15, 8189, 103}, {15,12284, 133}, {15,12285, 228}, {15,14332, 546},
{15,14333,1185}, {15,15356,1217}, {15,15357,2085}, {15,15868,2115}, {15,15869,2466},
{15,16124,3361}, {15,16125,3393}, {15,16252,3425}, {15,16253,3457}, {5, 1, 88},
};

```

Table 8. MPEG4 RVLC INTER VLC table.

Each entry is (code length, code word, and code pattern) stored in decimal.

For this example, take *pbits* equal to 7 which turns out to be the prefix size minimizing VLD table memory requirements. This means the left-most 7 bits of each codeword define which prefix-oriented table will include that codeword; if a codeword has less than 7 bits, then it will correspond to a range of prefixes. For example, the left column of following Table 9 shows the original MPEG4 RVLC Table entries (*length*, *pattern* (in decimal), *vlc_code* (in binary)), the middle column shows the corresponding partitioning into prefix-oriented tables for *pbits* = 7 with entries (*length-pbits*, *pattern* (in decimal), *vlc_code* & $(2^{\text{length-pbits}} - 1)$), and the right column has the entries of the middle column reorganized within each prefix-oriented table. In particular, Table 9 shows illustrates the following possibilities: no entries for *prefix* in various ranges, such as 0000000, ..., 0000011 (= 0x00, ..., 0x03 in hexadecimal); one entry for *prefix* in several ranges which arise when the codeword has less than *pbits* bits, such as *prefix* in the range 0x04-0x07 for codeword 00001; one entry for a single *prefix*, which arises for a codeword with exactly *pbits* bits, such as *prefix* = 0x1C for codeword 0011100; and multiple entries for fourteen different *prefix*, such *prefix* = 0x1E has two entries and *prefix* = 0x1F has 14 entries. The codeword entries in the center column of Table 9 are

in the format (*length-pbits*, *pattern*, prefix-stripped, left-shifted version of *vlc_code* in binary) and are listed in increasing order when interpreted as binary integers. After stripping off *prefix*, the maximum length of the codewords in a prefix-oriented table with entries in the center column is $plen_max = len_max - pbits$; for example, the prefix-oriented table with *prefix* = 0x1E has two entries and $plen_max = 1$ because after stripping off *prefix*, the two originally-8-bit codewords are now the two 1-bit codewords 0 and 1; similarly, the prefix-oriented table with *prefix* = 0x1F has 14 entries and $plen_max = 8$ because after stripping off *prefix*, the original codewords with lengths ranging from 9 to 15 bits now have 2 to 8 bits.

Original VLC Table (MPEG4 INTER RVLC)	Prefix-Oriented Table (<i>prefix</i> , <i>plen_max</i>)	Prefix-Oriented Table after reorganization with mask key
original table entries (N/A)	(<i>prefix</i> =0x0 - 0x3 N/A)	(N/A)
original table entries	(<i>prefix</i> =0x4 -0x7 <i>plen_max</i> =0)	(N/A)
5 88 00001-----	N/A	N/A
original table entries	(<i>prefix</i> =0x8 -0xf <i>plen_max</i> =0)	(N/A)
4 2 0001-----	N/A	N/A
original table entries	(<i>prefix</i> =0x10 -0x13 <i>plen_max</i> =0)	(N/A)
5 3 00100-----	N/A	N/A
original table entries	(<i>prefix</i> =0x14 -0x17 <i>plen_max</i> =0)	(N/A)
5 97 00101-----	N/A	N/A
original table entries	(<i>prefix</i> =0x18 -0x19 <i>plen_max</i> =0)	(N/A)
6 34 001100-----	N/A	N/A
original table entries	(<i>prefix</i> =0x1a -0x1b <i>plen_max</i> =0)	(N/A)
6 193 001101-----	N/A	N/A
original table entries	(<i>prefix</i> =0x1c -0x1c <i>plen_max</i> =0)	(N/A)
7 4 0011100-----	N/A	N/A
original table entries	(<i>prefix</i> =0x1d -0x1d <i>plen_max</i> =0)	(N/A)
7 66 0011101-----	N/A	N/A
original table entries	(<i>prefix</i> =0x1e <i>plen_max</i> =1)	(<i>prefix</i> =0x1e <i>plen_max</i> =1 mask key=0x0)
8 5 00111100-----	1 5 0	1 5 0
8 6 00111101-----	1 6 1	1 6 1

original table entries	(prefix=0x1f plen_max=8)	(prefix=0x1f plen_max=8 mask key=0xfc)
9 7 0011111100----- 9 36 0011111101----- 10 8 0011111100----- 10 9 0011111101----- 11 10 0011111100----- 11 11 0011111101----- 12 12 001111111100--- 12 39 001111111101--- 13 13 00111111111100-- 13 14 00111111111101-- 14 17 00111111111100- 14 18 00111111111101- 15 19 001111111111100 15 103 001111111111101	2 7 00----- 2 36 01----- 3 8 100----- 3 9 101----- 4 10 1100---- 4 11 1101---- 5 12 11100--- 5 39 11101--- 6 13 111100-- 6 14 111101-- 7 17 1111100- 7 18 1111101- 8 19 11111100 8 103 11111101	8 19 00000000 8 103 00000001 7 17 0000010- 7 18 0000011- 6 14 000010-- 6 13 000011-- 5 39 00010--- 5 12 00011--- 4 11 0010---- 4 10 0011---- 3 9 010----- 3 8 011----- 2 36 10----- 2 7 11-----
original table entries	(prefix=0x20 -0x23 plen max=0)	(N/A)
5 129 01000-----	N/A	N/A
original table entries	(prefix=0x24 -0x27 plen max=0)	(N/A)
5 161 01001-----	N/A	N/A
original table entries	(prefix=0x28 -0x29 plen max=0)	(N/A)
6 225 010100-----	N/A	N/A
original table entries	(prefix=0x2a -0x2b plen max=0)	(N/A)
6 257 010101-----	N/A	N/A
original table entries	(prefix=0x2c -0x2c plen max=0)	(N/A)
7 289 0101100-----	N/A	N/A
original table entries	(prefix=0x2d -0x2d plen max=0)	(N/A)
7 321 0101101-----	N/A	N/A
original table entries	(prefix=0x2e plen_max=1)	(prefix=0x2e plen_max=1 mask key=0x0)
8 35 01011100----- 8 98 01011101-----	1 35 0 1 98 1	1 35 0 1 98 1
original table entries	(prefix=0x2f plen_max=8)	(prefix=0x2f plen_max=8 mask key=0xfc)

9 67 010111100-----	2 67 00-----	8 133 00000000
9 162 010111101-----	2 162 01-----	8 228 00000001
10 37 010111100-----	3 37 100-----	7 41 0000010-
10 99 010111101-----	3 99 101-----	7 42 0000011-
11 38 0101111100----	4 38 1100----	6 16 000010--
11 68 0101111101----	4 68 1101----	6 15 000011--
12 69 01011111100---	5 69 11100---	5 100 00010---
12 100 010111111101---	5 100 11101---	5 69 00011---
13 15 0101111111100--	6 15 111100--	4 68 0010----
13 16 0101111111101--	6 16 111101--	4 38 0011----
14 41 01011111111100-	7 41 1111100-	3 99 010-----
14 42 01011111111101-	7 42 1111101-	3 37 011-----
15 133 010111111111100	8 133 11111100	2 162 10-----
15 228 010111111111101	8 228 11111101	2 67 11-----
original table entries	(prefix=0x30 -0x31 plen_max=0)	(N/A)
6 2145 011000-----	N/A	N/A
original table entries	(prefix=0x32 -0x33 plen_max=0)	(N/A)
6 2177 011001-----	N/A	N/A
original table entries	(prefix=0x34 -0x34 plen_max=0)	(N/A)
7 353 0110100-----	N/A	N/A
original table entries	(prefix=0x35 -0x35 plen_max=0)	(N/A)
7 2273 0110101-----	N/A	N/A
original table entries	(prefix=0x36 plen_max=1)	(prefix=0x36 plen_max=1 mask_key=0x0)
8 130 01101100-----	1 130 0	1 130 0
8 385 01101101-----	1 385 1	1 385 1
original table entries	(prefix=0x37 plen_max=8)	(prefix=0x37 plen_max=8 mask_key=0xfc)
9 481 011011100-----	2 481 00-----	8 546 00000000
9 513 011011101-----	2 513 01-----	8 1185 00000001
10 194 0110111100-----	3 194 100-----	7 70 0000010-
10 226 0110111101-----	3 226 101-----	7 71 0000011-
11 131 01101111100----	4 131 1100----	6 101 000010--
11 163 01101111101----	4 163 1101----	6 40 000011--
12 195 011011111100---	5 195 11100---	5 227 00010---
12 227 011011111101---	5 227 11101---	5 195 00011---
13 40 0110111111100--	6 40 111100--	4 163 0010----
13 101 0110111111101--	6 101 111101--	4 131 0011----
14 70 01101111111100-	7 70 1111100-	3 226 010-----
14 71 01101111111101-	7 71 1111101-	3 194 011-----
15 546 011011111111100	8 546 11111100	2 513 10-----
15 1185 011011111111101	8 1185 11111101	2 481 11-----
original table entries	(prefix=0x38 -0x38 plen_max=0)	(N/A)
7 2305 0111000-----	N/A	N/A
original table entries	(prefix=0x39 -0x39 plen_max=0)	(N/A)
7 2337 0111001-----	N/A	N/A

original table entries	(prefix=0x3a plen_max=1)	(prefix=0x3a plen_max=1 mask key=0x0)
8 417 01110100----- 8 449 01110101-----	1 417 0 1 449 1	1 417 0 1 449 1
original table entries	(prefix=0x3b plen_max=8)	(prefix=0x3b plen_max=8 mask key=0xfc)
9 545 011101100----- 9 2082 011101101----- 10 258 0111011100----- 10 290 0111011101----- 11 322 01110111100----- 11 673 01110111101----- 12 354 011101111100--- 12 865 011101111101--- 13 132 0111011111100-- 13 164 0111011111101-- 14 102 01110111111100- 14 196 01110111111101- 15 1217 011101111111100 15 2085 011101111111101	2 545 00----- 2 2082 01----- 3 258 100----- 3 290 101----- 4 322 1100----- 4 673 1101----- 5 354 11100--- 5 865 11101--- 6 132 111100-- 6 164 111101-- 7 102 1111100- 7 196 1111101- 8 1217 11111100 8 2085 11111101	8 1217 00000000 8 2085 00000001 7 102 0000010- 7 196 0000011- 6 164 000010-- 6 132 000011-- 5 865 00010--- 5 354 00011--- 4 673 0010---- 4 322 0011---- 3 290 010----- 3 258 011----- 2 2082 10----- 2 545 11-----
original table entries	(prefix=0x3c plen_max=1)	(prefix=0x3c plen_max=1 mask key=0x0)
8 2050 01111000----- 8 2433 01111001-----	1 2050 0 1 2433 1	1 2050 0 1 2433 1
original table entries	(prefix=0x3d plen_max=8)	(prefix=0x3d plen_max=8 mask key=0xfc)
9 2529 011110100----- 9 2561 011110101----- 10 577 0111101100----- 10 609 0111101101----- 11 705 01111011100----- 11 737 01111011101----- 12 897 011110111100--- 12 929 011110111101--- 13 259 0111101111100-- 13 386 0111101111101-- 14 291 01111011111100- 14 418 01111011111101- 15 2115 011110111111100 15 2466 011110111111101	2 2529 00----- 2 2561 01----- 3 577 100----- 3 609 101----- 4 705 1100----- 4 737 1101----- 5 897 11100--- 5 929 11101--- 6 259 111100-- 6 386 111101-- 7 291 1111100- 7 418 1111101- 8 2115 11111100 8 2466 11111101	8 2115 00000000 8 2466 00000001 7 291 0000010- 7 418 0000011- 6 386 000010-- 6 259 000011-- 5 929 00010--- 5 897 00011--- 4 737 0010---- 4 705 0011---- 3 609 010----- 3 577 011----- 2 2561 10----- 2 2529 11-----
original table entries	(prefix=0x3e plen_max=8)	(prefix=0x3e plen_max=8 mask key=0xfc)

9 2593 011111000----- 9 2625 011111001----- 10 641 0111110100----- 10 2721 0111110101----- 11 769 01111101100----- 11 801 01111101101----- 12 2083 011111011100--- 12 2146 011111011101--- 13 961 0111110111100-- 13 993 0111110111101-- 14 450 01111101111100- 14 482 01111101111101- 15 3361 011111011111100 15 3393 011111011111101	2 2593 00----- 2 2625 01----- 3 641 100----- 3 2721 101----- 4 769 1100----- 4 801 1101----- 5 2083 11100--- 5 2146 11101--- 6 961 111100-- 6 993 111101-- 7 450 1111100- 7 482 1111101- 8 3361 11111100 8 3393 11111101	8 3361 00000000 8 3393 00000001 7 450 0000010- 7 482 0000011- 6 993 000010-- 6 961 000011-- 5 2146 00010--- 5 2083 00011--- 4 801 0010---- 4 769 0011---- 3 2721 010----- 3 641 011----- 2 2625 10----- 2 2593 11-----
original table entries	(prefix=0x3f plen_max=8)	(prefix=0x3f plen_max=8 mask key=0x7c)
10 2753 0111111000----- 10 2785 0111111001----- 11 833 01111110100----- 11 2051 01111110101----- 12 2178 011111101100--- 12 2977 011111101101--- 13 1025 0111111011100-- 13 1057 0111111011101-- 14 514 01111110111100- 14 1089 01111110111101- 15 3425 011111101111100 15 3457 011111101111101 11 2114 01111111000----- 11 2881 01111111001----- 12 3009 011111110100--- 12 3041 011111110101--- 13 2052 0111111101100-- 13 2210 0111111101101-- 14 1121 01111111011100- 14 1153 01111111011101- 12 3073 011111111000--- 12 3105 011111111001--- 13 2242 0111111110100-- 13 2274 0111111110101-- 14 2053 01111111101100- 14 2084 01111111101101- 13 2306 0111111111000-- 13 2338 0111111111001-- 14 2370 01111111110100- 14 2402 01111111110101- 14 2434 01111111111000- 14 3265 01111111111001-	3 2753 000----- 3 2785 001----- 4 833 0100----- 4 2051 0101----- 5 2178 01100--- 5 2977 01101--- 6 1025 011100-- 6 1057 011101-- 7 514 0111100- 7 1089 0111101- 8 3425 01111100 8 3457 01111101 4 2114 1000----- 4 2881 1001----- 5 3009 10100--- 5 3041 10101--- 6 2052 101100-- 6 2210 101101-- 7 1121 1011100- 7 1153 1011101- 5 3073 11000--- 5 3105 11001--- 6 2242 110100-- 6 2274 110101-- 7 2053 1101100- 7 2084 1101101- 6 2306 111000-- 6 2338 111001-- 7 2370 1110100- 7 2402 1110101- 7 2434 1111000- 7 3265 1111001-	8 3425 00000000 8 3457 00000001 7 514 0000010- 7 1089 0000011- 6 1057 000010-- 6 1025 000011-- 5 2977 00010--- 5 2178 00011--- 4 2051 0010---- 4 833 0011---- 3 2785 010----- 3 2753 011----- 7 2434 1000110- 7 3265 1000111- 7 2370 1001010- 7 2402 1001011- 6 2338 100110-- 6 2306 100111-- 7 2053 1010010- 7 2084 1010011- 6 2274 101010-- 6 2242 101011-- 5 3105 10110--- 5 3073 10111--- 7 1121 1100010- 7 1153 1100011- 6 2210 110010-- 6 2052 110011-- 5 3041 11010--- 5 3009 11011--- 4 2881 1110---- 4 2114 1111----
original table entries	(prefix=0x40 plen_max=7)	(prefix=0x40 plen_max=7 mask key=0x2)

14 3297 100000000000010-	7 3297 0000010	7 3297 0000000
14 3329 100000000000011-	7 3329 0000011	7 3329 0000001
13 3201 10000000000010--	6 3201 000010-	6 3233 000010-
13 3233 10000000000011--	6 3233 000011-	6 3201 000011-
12 3137 1000000000010---	5 3137 00010--	5 3137 00010--
12 3169 1000000000011---	5 3169 00011--	5 3169 00011--
11 2913 100000000010----	4 2913 0010---	4 2913 0010---
11 2945 100000000011----	4 2945 0011---	4 2945 0011---
10 2817 10000000010-----	3 2817 010----	3 2817 010----
10 2849 10000000011-----	3 2849 011----	3 2849 011----
9 2657 1000000010-----	2 2657 10-----	2 2657 10-----
9 2689 1000000011-----	2 2689 11-----	2 2689 11-----
original table entries	(prefix=0x41 plen_max=1)	(prefix=0x41 plen_max=1 mask key=0x0)
8 2465 10000010-----	1 2465 0	1 2465 0
8 2497 10000011-----	1 2497 1	1 2497 1
original table entries	(prefix=0x42 plen_max=0)	(N/A)
7 2369 1000010-----	N/A	N/A
original table entries	(prefix=0x43 plen_max=0)	(N/A)
7 2401 1000011-----	N/A	N/A
original table entries	(prefix=0x44 plen_max=0)	(N/A)
6 2209 100010-----	N/A	N/A
original table entries	(prefix=0x46 plen_max=0)	(N/A)
6 2241 100011-----	N/A	N/A
original table entries	(prefix=0x48 plen_max=0)	(N/A)
5 2081 10010-----	N/A	N/A
original table entries	(prefix=0x4c plen_max=0)	(N/A)
5 2113 10011-----	N/A	N/A
original table entries	(prefix=0x50 plen_max=0)	(N/A)
4 65 1010-----	N/A	N/A
original table entries	(prefix=0x58 plen_max=0)	(N/A)
4 2049 1011-----	N/A	N/A
original table entries	(prefix=0x60 plen_max=0)	(N/A)
3 1 110-----	N/A	N/A
original table entries	(prefix=0x70 plen_max=0)	(N/A)
3 33 111-----	N/A	N/A

Table 9. The VLC Table in Table 8 decomposed into prefix-oriented tables

Note that one effect of partitioning into prefix-oriented tables and then into sub-tables for the AUVLD control table and code table construction as compared

to a direct VLD control table and code table construction is the decrease in the number of repetitions in the code table due to a smaller range of codeword lengths in the sub-tables. For example, dividing Table 8 directly into sub-tables would mean that all codewords beginning with 01 would be in the same sub-table and these codewords vary in length from 5 bits to 15 bits. In contrast, in the prefix-oriented tables with prefixes = 01xxxxx, the prefix-stripped codewords varying in length from 2 bits to 8 bits in five of the prefix-oriented tables and from 4 bits to 8 bits in the largest prefix-oriented table which has prefix 0111111; see Table 9.

Next, for each prefix-oriented table shown in the center column of Table 9, consider all possible bit masks, *mask_key*, of *plen_max* bits, and XOR (bit-by-bit) the left-shifted, *prefix*-stripped codewords with such *mask_key*. Then reorganize the XORED prefix-oriented table with increasing order by interpreting as binary integers; see the third column of Table 9. For example, the prefix-oriented table for *prefix* = 0x1F has *plen_max* = 8, and thus with 8-bit *mask_key* = 0xFC (=1111 1100) the *prefix*-stripped, left-shifted codewords have their first 6 bits reversed by the XORing. These XORED codewords are reorganized as illustrated in the right column of Table 9 which shows the entry 8 19 11111100 of the center column becoming entry 8 19 00000000 of the right column and thus moving in the ordering from second largest in the center column to smallest in the right column. Each possible 8-bit *mask_key* is used, and the one yielding the smallest memory requirements for the resulting control and code tables (next paragraph) will be selected. Very roughly, the *mask_key* which makes the longest codewords in the prefix-oriented table have the most leading 0s will minimize memory requirements when using the 0-leading construction of VLD control and code tables of the foregoing sections. Indeed, the *reverse* used for 1-leading codes in the foregoing to convert them to 0-leading codes is the special case of taking *mask_key* = 111...11. The pair (*plen_max*, *mask_key*) is the entry in AUVLD prefix table (VLDPrefixTab in Figure 1b) corresponding to *prefix*; this prefix table has 2^{pbits} entries.

For a prefix-oriented table with none or only one entry (so *plen_max* = 0), construct the corresponding AUVLD control table with only one entry. Of course, when *plen_max* = 0, the first bits of *prefix* are the codeword.

For each prefix-oriented table with multiple entries (XORed *prefix*-stripped, left-shifted codewords), construct a control table plus a code table in the same manner as described for the 0-leading construction of the UVLD code table and UVLD control table in the foregoing sections. The overall AUVLD control table is two-dimensional in the sense that the subtables of the AUVLD code table are indexed by both *prefix* and *subtab_id*, the same subtable index as in a UVLD control table. However, following section 6 will merge the prefix table into the control table to reduce memory, and Table 12 shows the ultimate VLD code table and VLD control table for Table 8.

5. Advanced universal VLD decoding

As illustrated in Figure 1b, AUVLD decodes using the tables constructed in the foregoing section (the prefix table *VLDPrefixTab*[], the two-dimensional VLD control table *VLDCtlTab*[][.], and the VLD code table *VLDCodeTab*[]) including the corresponding parameters *len_max* and *pbits*. Indeed, given the following data types:

```
typedef struct vldcodetab{
    char      length;
    short     pattern
} VLDCodeTab;

typedef struct vldctltab{
    char      shift;
    short     offset;
} VLDCtlTab;

typedef struct vldprefixtab{
    char      plen_max;
    int       mask_key;
} VLDPrefixTab;
```

then AUVLD decodes as follows:

1. Get *prefix* for the next codeword (which encodes the next symbol) from the first *pbits* of the bitstream starting at the current decoding position.

2. Obtain the maximum code length (*plen_max*) and mask (*mask_key*) for the prefix-oriented table corresponding to *prefix* by look up in the VLD prefix table:

```
plen_max = VLDPrefixTab[prefix].plen_max;
mask_key = VLDPrefixTab[prefix].mask_key;.
```

3. Get *value*, the next *plen_max* bits (following the *pbits* bits already read) from the bitstream; note that *plen_max* may be 0.

4. Perform XOR with the mask key from step 2: *value* = XOR(*value*, *mask_key*) where *mask_key* was found in step 2 as part of the look up.

5. Use the XORed *value* from step 4 to determine the sub-table index, *subtab_id*, for addressing the VLD control table: if *value* is 0 (e.g., *plen_max* = 0), then *subtab_id* = 0; else *subtab_id* = (int) $\log_2(\text{value}) + 1$.

6. Get *shift* and *offset* from the (two-dimensional) VLD control table; *prefix* is the first index and *subtab_id* the second. In Figure 1b *prefix* controls SWITCH to set VLD control table first index.

```
shift = VLDCtlTab[prefix][subtab_id].shift;
offset = VLDCtlTab[prefix][subtab_id].offset;
```

7. Compute the *index* to use in the VLD code table from *shift* and *offset*:
index = (*value* >> *shift*) + *offset*; this is the same as with UVLD.

```
index = (value >> shift) + offset;
```

8. Use *index* to look up *length* and *pattern* in the VLD code table; this is the same as with UVLD.

```
length = VLDCodeTab[index].length;
pattern = VLDCodeTab[index].pattern;
```

9. Set the next decoder position by moving the current decoding position by *length* bits; and interpret *pattern* to decode the symbol; again the same as with UVLD.

10. Loop back to step 1 for decoding the next symbol.

Note that errors are detected when the bits read do not correspond to a codeword and thus do not lead to an entry in the VLD code table.

The following table 10 shows pseudocode for the decoding of Figure 1b.

```

int AdvancedUniversalVLD(
    Bitstream *stream,          /* pointer of bitstream
*/
    VLDCodeTab *vldtab,        /* pointer of VLD code table
*/
    VLDCtlTab **vldctl,        /* pointer of VLD control table
*/
    int len_max,                /* maximum code length in the VLC table
*/
    int pbits,                  /* number of prefix bits
*/
    VLDPrefixTab *vldprefixtab, /* pointer of VLD prefix table
*/
    char *err_flag)            /* err_flag =1 ->error detected, err_flag=0->decoding OK
{
    int value, subtab_id, index;
    int prefix, mask_key, plen_max;
    /*=====
*/
    /* get the value of next "len_max" bits in the bitstream
    */
    /*=====
*/
    value = next_bits(stream, len_max);

    /*=====*/
    /* get the value of prefix
    */
    /*=====
*/
    prefix = value>>(len_max-pbits);

    /*=====
*/
    /* get maximum code length and mask key of prefix-oriented table
    */
    /*=====
*/
    plen_max = vldprefixtab[prefix].plen_max;
    mask_key = vldprefixtab[prefix].mask_key;

    /*=====*/
    /* get the value of next "plen_max" bits
    */
    /*=====
*/
    value = (value>>(len_max-pbits-plen_max)) & ((1<<plen_max)-1);

    /*=====*/

```

```

    /* Perform XOR operation
*/
/*=====
= */
    value = value ^ mask_key;

/*=====
= */
    /* determine the sub-table index according to the value. TMS320C6X and TMS320C54X have
    */
    /* special instructions for such an operation
    */
    /*=====
= */
    if (value==0) subtab_id = 0; else subtab_id = (int) log2(value) + 1;

/*=====
= */
    /* get index in the VLD code table
    */
    /*=====
= */
    index = vldctl[prefix][subtab_id].offset + (value>>vldctl[prefix][subtab_id].shift);

/*=====
= */
    /* decide if an decoding error is detected
    */
    /*=====
= */
    if (vldtab[index].length ==0) *err_flag =1; else *err_flg=0;

/*=====
= */
    /* update the current decoding position in the bitstream
    */
    /*=====
= */
    if (*err_flg==0) flush_bits(stream, vldtab[index].length);

/*=====
= */
    /* return the decoded coding pattern
    */
    /*=====
= */
    return vldtab[index].pattern;
}

```

Table 10, pseudocode for the AUVLD of Figure 1b

6. AUVLD efficient table storage

The AUVLD of the preceding sections shares the same table construction as the UVLD; the only difference is the UVLD table is stored in a more efficient way. According, a variation of AUVLD provides a more efficient storage using a another table format as follows. First, note that the prefix table has 2^{pbits} entries and that each entry could be 5 bytes: 1 byte for *plen_max* and 4 bytes for the *mask_key*. But the majority of prefix-oriented tables have none or only one entry and do not use or need *mask_key* stored. In the Table 9 MPEG-4 RVLC example, only 14 of the 128 prefixes have a prefix-oriented table with multiple entries. And storing the control table as a two-dimensional array wastes memory in that the number of entries is the product of 2^{pbits} times $\max\{plen_max_j + 1\}$ over $j = 0, 1, \dots, 2^{pbits} - 1$. Thus, a prefix-oriented table with only one entry has to maintain 9 entries to make the two-dimensional control table addressing (*prefix*, *subtab_id*) possible. The two-dimensional control table in this case will have a total of 128x9 entries, most of the entries are dummy entries. Therefore, second preferred embodiments AUVLD combine the prefix table and the two-dimensional control table into a one-dimensional control table, VLDCtlTab[,] as illustrated in Figure 1c. The VLD code table remains unchanged, and the combination of the prefix and control tables is as follows.

1. combine each entry (*shift*, *offset*) of the control table of each prefix-oriented table into a 16-bit word, with 5-bit *shift* followed by 11-bit *offset*.
2. embed *mask_key* of each prefix-oriented table into its control table VLDCtlTab by applying the following rules where VLDCtlTab is the VLD control table, *plen_max* and *mask_key* are the maximum code length and mask key for the prefix-oriented table, and *len_max* is the maximum code length for the VLC table :
 - a) if *plen_max* = 0, then no *mask_key* is embedded;
 - b) if *plen_max* > 0 and *len_max* ≤ 16, then put *mask_key* in VLDCtlTab[−1];
 - c) if *plen_max* > 0 and *len_max* > 16, then put (*mask_key* >> 16) in VLDCtlTab[−2] and (*mask_key* & 0xFFFF) in VLDCtlTab[−1]
3. cascade the individual VLD control tables (with *mask_key* embedded) together into a lined-up one-dimensional VLD control table, in the order of

increasing prefix value. Dummy VLD control table entries are eliminated in this step.

4. place 2^{pbits} 16-bit entries in front of the lined-up one-dimensional VLD control table to make up a combined VLD control table; see Figure 1c. These first 2^{pbits} 16-bit entries are addressed by *prefix*. Each of these entries is a doublet (*plen_max*, *ptab_id*) where *plen_max* is the maximum code length within the a prefix-oriented table and *ptab_id* points to the start index of the VLD control table of a prefix-oriented table in the combined one-dimensional control table; see Figure 1c. The entry (*plen_max*, *ptab_id*) is combined into 16 bits by allocation of 5 bits for *plen_max* followed by 11 bits for *ptab_id*.

5. combine each entry (*length*, *pattern*) in the VLD code table into 16 bits with (16 – *symbol_bits*)-bit *length* followed by (*symbol_bits*)-bit *pattern*. where *symbol_bits* is the number of bits used in storing *pattern*.

The alternative embodiment AUVLD decoder uses these two tables together with the parameters *len_max*, *pbits*, and *symbol_bits* for decoding. Note that all possible tables corresponding to combinations values of *pbits* and the various *mask_keys* are computed, and the values leading to the smallest memory required are selected. Of course, the universal decoding would also work for the tables derived for other values of *pbits* and the various *mask_keys*; only the proper table format is needed for this particular AUVLD.

7. AUVLD decoding with efficient storage

As illustrated in Figure 1c, with the preceding section efficient storage of merged prefix and control tables, AUVLD decodes using the VLD control table, VLD code table, and parameters *len_max*, *pbits*, and *symbol_bits* as follows:

1. Get *prefix* for the next codeword by reading the first *pbits* from the bitstream starting from the current decoding position.
2. Obtain the maximum code length (*plen_max*) and prefix-oriented table index (*ptab_id*) for the prefix-oriented table corresponding to *prefix* by look up in the one-dimensional control table using $plen_max = \text{VLDCtlTab}[prefix] \gg 11$ (the

first 5 bits of the table entry) and $ptab_id = VLDCtlTab[prefix] \& 0x7FF$ (the last 11 bits).

3. Get $mask_key$ by

if $plen_max > 0$ and $len_max \leq 16$, then $mask_key = VLDCtlTab[ptab_id-1]$

if $len_max > 0$ and $len_max > 16$, then $mask_key = (VLDCtlTab[ptab_id-2] \ll 16) + VLDCtlTab[ptab_id-1]$

4. Get, $value$, the next $plen_max$ bits (following the $pbits$ bits already read) from the bitstream; note that $plen_max$ may be 0.

5. Perform XOR: $value = XOR(value, mask_key)$ where $mask_key$ was found in step 3 look up.

6. Use $value$ from step 5 to determine the sub-table index, $subtab_id$, for addressing the VLD control table: if $value$ is 0 (e.g., $plen_max = 0$), then $subtab_id = 0$; else $subtab_id = (int) \log_2(value) + 1$.

7. Get $shift$ and $offset$ from the (one-dimensional) VLD control table by $shift = VLDCtlTab[ptab_id+subtab_id] \gg 11$ (first 5 bits of entry) and $offset = VLDCtlTab[ptab_id+subtab_id] \& 0x7FF$ (last 11 bits)

8. Compute the $index$ in the VLD code table from $shift$ and $offset$: $index = (value \gg shift) + offset$; this is the same as with UVLD.

9. Use $index$ to look up $length$ and $pattern$ for the current encoded symbol in the VLD code table by $length = VLDCodeTab[index] \gg symbol_bits$ (first bits of entry) and $pattern = VLDCodeTab[index] \& 2^{symbol_bits}-1$ (last bits) .

10. Set the next decoder position by moving the current decoding position by $length$ bits; and interpret $pattern$ to get the decoded symbol; again the same as with UVLD.

11. Loop to step 1 for the next symbol decoding.

Following Table 11 shows pseudocode for the decoding of Figure 1c, and Table 12 shows the corresponding VLD code table and VLD control table constructed from Table 8. Note that the VLD code table has 295 entries and the VLD control table has 256 entries where each entry is one 2-byte word; whereas, the original MPEG-4 RVLC table has 170 entries with each entry ($length$, $pattern$, vlc_code).

Thus the preferred embodiment conversion to universal decoding format does not greatly expand the VLC table.

```

int AdvancedUniversalVLD(
    Bitstream *stream,          /* pointer of bitstream
*/
    Unsigned short *vldtab,     /* pointer of VLD code table
*/
    Unsigned short *vldctl,     /* pointer of VLD control table
*/
    int len_max,               /* maximum code length in the VLD table
*/
    int pbits,                 /* number of prefix bits
*/
    int symbol_bits,           /* number of bits used for saving pattern in VLD code table entry
*/
    char *err_flag)            /* err_flag = 1 ->error detected, err_flag=0->decoding OK
*/
{
    int value, subtab_id, index;
    int prefix, mask_key, plen_max;
    /*=====
*/
    /* get the value of next "len_max" bits in the bitstream
    */
    /*=====
*/
    value = next_bits(stream, len_max);

    /*=====
*/
    /* get the value of prefix
    */
    /*=====
*/
    prefix = value>>(len_max-pbits);

    /*=====
*/
    /* get maximum code length and mask key of prefix-oriented table
    */
    /*=====
*/
    plen_max = vldctl[prefix]>>11;
    ptab_id = vldctl[prefix]&0x7ff;
    if (len_max ≤ 16 && plen_max > 0) mask_key = vldctl[ptab_id-1];
    else if (len_max > 16 && plen_max > 0)
        mask_key = (vldctltab [ptab_id-2]<<16) + vldctl[ptab_id-1];
    else mask_key = 0;

    /*=====
*/

```



```

    /* get the value of next "plen_max" bits
    */
    /*=====
    */
    value = (value>>(len_max-pbits-plen_max)) & ((1<<plen_max)-1);

    /*=====*/
    /* Perform XOR operation
    */
    /*=====
    */
    value = value ^ mask_key

    /*=====
    */
    /* determine the sub-table index according to the value. TMS320C6X and TMS320C54X have
    */
    /* special instructions for such an operation
    */
    /*=====
    */
    if (value==0) subtab_id = 0; else subtab_id = (int) log2(value) + 1;

    /*=====
    */
    /* get index in the VLD code table
    */
    /*=====
    */
    index = vldctl[ptab_id+subtab_id]&0x7ff + (value>> (vldctl[ptab_id+subtab_id]>>11));

    /*=====
    */
    /* decide if an decoding error is detected
    */
    /*=====
    */
    if ((vldtab[index]>>symbol_bits) ==0) *err_flag =1; else *err_flag=0;

    /*=====
    */
    /* update the current decoding position in the bitstream
    */
    /*=====
    */
    if (*err_flag==0) flush_bits(stream, vldtab[index]>>symbol_bits));

    /*=====
    */
    /* return the decoded coding pattern
    */
    /*=====
    */
    return vldtab[index]&(2symbol_bits-1);
}

```

Table 11, pseudocode for the AUVLD of Figure 1c

```

/* 4 bit len + 12 bit pattern (symbol_bits=12), MPEG4_RVLC_INTER */
static unsigned short
MPEG4_RVLC_INTER_dcdtab[295]={
0x0000, 0x5058, 0x5058, 0x5058, 0x5058, 0x4002, 0x4002, 0x4002,
0x4002, 0x4002, 0x4002, 0x4002, 0x4002, 0x5003, 0x5003, 0x5003,
0x5003, 0x5061, 0x5061, 0x5061, 0x5061, 0x6022, 0x6022, 0x60c1,
0x60c1, 0x7004, 0x7042, 0x8005, 0x8006, 0xf013, 0xf067, 0xe011,
0xe012, 0xd00e, 0xd00d, 0xc027, 0xc00c, 0xb00b, 0xb00a, 0xa009,
0xa008, 0x9024, 0x9007, 0x5081, 0x5081, 0x5081, 0x5081, 0x50a1,
0x50a1, 0x50a1, 0x50a1, 0x60e1, 0x60e1, 0x6101, 0x6101, 0x7121,
0x7141, 0x8023, 0x8062, 0xf085, 0xf0e4, 0xe029, 0xe02a, 0xd010,
0xd00f, 0xc064, 0xc045, 0xb044, 0xb026, 0xa063, 0xa025, 0x90a2,
0x9043, 0x6861, 0x6861, 0x6881, 0x6881, 0x7161, 0x78e1, 0x8082,
0x8181, 0xf222, 0xf4a1, 0xe046, 0xe047, 0xd065, 0xd028, 0xc0e3,
0xc0c3, 0xb0a3, 0xb083, 0xa0e2, 0xa0c2, 0x9201, 0x91e1, 0x7901,
0x7921, 0x81a1, 0x81c1, 0xf4c1, 0xf825, 0xe066, 0xe0c4, 0xd0a4,
0xd084, 0xc361, 0xc162, 0xb2a1, 0xb142, 0xa122, 0xa102, 0x9822,
0x9221, 0x8802, 0x8981, 0xf843, 0xf9a2, 0xe123, 0xe1a2, 0xd182,
0xd103, 0xc3a1, 0xc381, 0xb2e1, 0xb2c1, 0xa261, 0xa241, 0x9a01,
0x99e1, 0xfd21, 0xfd41, 0xe1c2, 0xe1e2, 0xd3e1, 0xd3c1, 0xc862,
0xc823, 0xb321, 0xb301, 0xaaa1, 0xa281, 0x9a41, 0x9a21, 0xfd61,
0xfd81, 0xe202, 0xe441, 0xd421, 0xd401, 0xcba1, 0xc882, 0xb803,
0xb341, 0xaae1, 0xaac1, 0x0000, 0x0000, 0x0000, 0x0000, 0x0000,
0x0000, 0xe982, 0xeccl, 0x0000, 0x0000, 0x0000, 0xe942, 0xe962, 0xd922,
0xd922, 0xd902, 0xd902, 0x0000, 0x0000, 0xe805, 0xe824, 0xd8e2,
0xd8e2, 0xd8c2, 0xd8c2, 0xcc21, 0xcc21, 0xcc21, 0xcc21, 0xcc01,
0xcc01, 0xcc01, 0xcc01, 0x0000, 0x0000, 0xe461, 0xe481, 0xd8a2,
0xd8a2, 0xd804, 0xd804, 0xcbe1, 0xcbe1, 0xcbe1, 0xcbe1, 0xcbe1,
0xcbe1, 0xcbe1, 0xcbe1, 0xbb41, 0xbb41, 0xbb41, 0xbb41, 0xbb41,
0xbb41, 0xbb41, 0xbb41, 0xb842, 0xb842, 0xb842, 0xb842, 0xb842,
0xb842, 0xb842, 0xb842, 0xece1, 0xed01, 0xdca1, 0xdc81, 0xcc41,
0xcc61, 0xbb61, 0xbb81, 0xab01, 0xab21, 0x9a61, 0x9a81, 0x89a1,
0x89c1, 0x7941, 0x7961, 0x68a1, 0x68a1, 0x68c1, 0x68c1, 0x5821,
0x5821, 0x5821, 0x5821, 0x5841, 0x5841, 0x5841, 0x5841, 0x4041,
0x4041, 0x4041, 0x4041, 0x4041, 0x4041, 0x4041, 0x4041, 0x4801,
0x4801, 0x4801, 0x4801, 0x4801, 0x4801, 0x4801, 0x4801, 0x3001,
0x3001, 0x3001, 0x3001, 0x3001, 0x3001, 0x3001, 0x3001, 0x3001,
0x3001, 0x3001, 0x3001, 0x3001, 0x3001, 0x3001, 0x3001, 0x3021,
0x3021, 0x3021, 0x3021, 0x3021, 0x3021, 0x3021, 0x3021, 0x3021,
0x3021, 0x3021, 0x3021, 0x3021, 0x3021, 0x3021, 0x3021, 0x3021,
};

/* prefix_bits = 7, len_max = 15 */
/* 5 bit SHIFT + 11 bit OFFSET */
static unsigned short
MPEG4_RVLC_dcdctl[256]={
0x0080, 0x0080, 0x0080, 0x0080, 0x0081, 0x0081, 0x0081, 0x0081,
0x0082, 0x0082, 0x0082, 0x0082, 0x0082, 0x0082, 0x0082, 0x0082,
0x0083, 0x0083, 0x0083, 0x0083, 0x0084, 0x0084, 0x0084, 0x0084,
0x0085, 0x0085, 0x0086, 0x0086, 0x0087, 0x0088, 0x008a, 0x408d,
0x0096, 0x0096, 0x0096, 0x0096, 0x0097, 0x0097, 0x0097, 0x0097,

```

```

0x0098, 0x0098, 0x0099, 0x0099, 0x009a, 0x009b, 0x089d, 0x40a0,
0x00a9, 0x00a9, 0x00aa, 0x00aa, 0x00ab, 0x00ac, 0x08ae, 0x40b1,
0x00ba, 0x00bb, 0x08bd, 0x40c0, 0x08ca, 0x40cd, 0x40d7, 0x40e1,
0x38eb, 0x08f4, 0x00f6, 0x00f7, 0x00f8, 0x00f8, 0x00f9, 0x00f9,
0x00fa, 0x00fa, 0x00fa, 0x00fa, 0x00fb, 0x00fb, 0x00fb, 0x00fb,
0x00fc, 0x00fc, 0x00fc, 0x00fc, 0x00fc, 0x00fc, 0x00fc, 0x00fc,
0x00fd, 0x00fd, 0x00fd, 0x00fd, 0x00fd, 0x00fd, 0x00fd, 0x00fd,
0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe,
0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe, 0x00fe,
0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff,
0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff, 0x00ff,
0x0000, 0x0001, 0x0005, 0x000d, 0x0011, 0x0015, 0x0017, 0x0019,
0x001a, 0x0000, 0x001b, 0x001b, 0x00fc, 0x001d, 0x001d, 0x401f,
0x081d, 0x101f, 0x1821, 0x2023, 0x2825, 0x3027, 0x002b, 0x002f,
0x0033, 0x0035, 0x0037, 0x0038, 0x0000, 0x0039, 0x0039, 0x00fc,
0x003b, 0x003b, 0x403d, 0x083b, 0x103d, 0x183f, 0x2041, 0x2843,
0x3045, 0x0049, 0x004b, 0x004d, 0x004e, 0x0000, 0x004f, 0x004f,
0x00fc, 0x0051, 0x0051, 0x4053, 0x0851, 0x1053, 0x1855, 0x2057,
0x2859, 0x305b, 0x005f, 0x0060, 0x0000, 0x0061, 0x0061, 0x00fc,
0x0063, 0x0063, 0x4065, 0x0863, 0x1065, 0x1867, 0x2069, 0x286b,
0x306d, 0x0000, 0x0071, 0x0071, 0x00fc, 0x0073, 0x0073, 0x4075,
0x0873, 0x1075, 0x1877, 0x2079, 0x287b, 0x307d, 0x00fc, 0x0081,
0x0081, 0x4083, 0x0881, 0x1083, 0x1885, 0x2087, 0x2889, 0x308b,
0x007c, 0x008f, 0x008f, 0x4091, 0x088f, 0x1091, 0x1893, 0x2095,
0x2897, 0x085b, 0x0002, 0x00db, 0x00db, 0x38dd, 0x08db, 0x10dd,
0x18df, 0x20e1, 0x28e3, 0x0000, 0x00e7, 0x00e7, 0x00e9, 0x00ea,
0x00eb, 0x00ed, 0x00ef, 0x00f3, 0x00f7, 0x00ff, 0x0107, 0x0117,

```

```
};
```

Table 12. VLD Code table and control table for MPEG4 INTER RVLC Table 8

8. Modifications

Various modifications can be made to the preferred embodiments while retaining the features of universal variable length decoding (UVLD) using a translation of a given VLC table into a VLD control table plus VLD code table and a prefix table (combinable with the control table) for Advanced UVLD. In particular, the exemplary VLC tables used could be replaced by others with differing bit allocations and symbol bits, bit complementary versions could be used, reordering the VLD code table entries and corresponding recomputation of the index, a three-tier or more decomposition (e.g., two-tier prefixes) for very large VLC tables may reduce memory required, and so forth.